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MINIMUM RANKS OF SIGN PATTERNS VIA SIGN VECTORS AND DUALITY

MARINA ARAV, FRANK J. HALL, ZHONGSHAN LI, HEIN VAN DER HOLST, JOHN SINKOVIC, AND LIHUA ZHANG

Abstract. A sign pattern matrix is a matrix whose entries are from the set \{+,-,0\}. The minimum rank of a sign pattern matrix \(A\) is the minimum of the ranks of the real matrices whose entries have signs equal to the corresponding entries of \(A\). It is shown in this paper that for any \(m \times n\) sign pattern \(A\) with minimum rank \(n-2\), rational realization of the minimum rank is possible. This is done using a new approach involving sign vectors and duality. It is shown that for each integer \(n \geq 9\), there exists a nonnegative integer \(m\) such that there exists an \(m \times n\) sign pattern matrix with minimum rank \(n-3\) for which rational realization is not possible. A characterization of \(m \times n\) sign patterns \(A\) with minimum rank \(n-1\) is given (which solves an open problem in Brualdi et al. [R. Brualdi, S. Fallat, L. Hogben, B. Shader, and P. van den Driessche. Final report: Workshop on Theory and Applications of Matrices Described by Patterns. Banff International Research Station, Jan. 31 – Feb. 5, 2010.], along with a more general description of sign patterns with minimum rank \(r\), in terms of sign vectors of certain subspaces. Several related open problems are stated along the way.

Key words. Sign pattern matrix, Sign vectors, Minimum rank, Maximum rank, Rational realization, Oriented matroid duality.

AMS subject classifications. 15B35, 15B36, 52C40.

1. Introduction. An important part of combinatorial matrix theory is the study of sign pattern matrices, which has been the focus of extensive research for the last 50 years ([8], [15]). A sign pattern matrix is a matrix whose entries are from the set \{+, -, 0\}. For a real matrix \(B\), \text{sign}(B)\) is the sign pattern matrix obtained by replacing each positive (respectively, negative, zero) entry of \(B\) by + (respectively, - , 0). For a sign pattern matrix \(A\), the sign pattern class (also known as the qualitative class) of \(A\), denoted \(Q(A)\), is defined as

\[
Q(A) = \{ B : B \text{ is a real matrix and } \text{sign}(B) = A \}.
\]

The minimum rank of a sign pattern matrix \(A\), denoted \(mr(A)\), is the minimum of the ranks of the real matrices in \(Q(A)\). Determination of the minimum rank of a sign

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minimum rank of a sign pattern matrix in general is a longstanding open problem (see [17]) in combinatorial matrix theory. Recently, there has been a number of papers concerning this topic, for example [1]–[12], [14]–[22]. In particular, matrices realizing the minimum rank of a sign pattern have applications in the study of neural networks [10] and communication complexity [22].

More specifically, in the study of communication complexity in computer science, \((+,-)\) sign pattern matrices arise naturally; the minimum rank (also known as the sign-rank) of the sign patterns plays an important role, as the minimum rank of a corresponding sign pattern matrix essentially determines the unbounded-error randomized communication complexity of a function.

In mathematics, rational realization is an important theme. For example, the study of the existence of rational (or integer) solutions of Diophantine equations is well-known. Steinitz’s celebrated theorem stating that the combinatorial type of every 3-polytope is rationally realizable is a far-reaching result [27]. In combinatorics, rational realizability of certain point-line configurations is an important problem ([13, 25]); this is closely related to the realizability of the minimum rank of a sign pattern with minimum rank 3 [18]. More generally, the rational realizability of the minimum rank of a sign pattern matrix is equivalent to the rational realizability of a certain point-hyperplane configuration.

In [1], several classes of sign patterns \(A\) for which rational realization of the minimum rank is guaranteed are identified, such as when every entry of \(A\) is nonzero, or the minimum rank of \(A\) is at most 2, or the minimum rank of \(A\) is at least \(n - 1\), where \(A\) is \(m \times n\). It has been shown in [19], through the use of a result in projective geometry, that rational realization of the minimum rank is not always possible. Specifically, in [19], the authors showed that there exists a \(12 \times 12\) sign pattern matrix with minimum rank 3 but there is no rational realization of rank 3 within the qualitative class of the sign pattern. Independently, Berman et al. [5] also provided an example of a sign pattern for which the rational minimum rank is strictly greater than the minimum rank over the reals. Both of these papers use techniques based on matroids. More recently, Jing et al. [18] found a \(9 \times 9\) sign pattern matrix with minimum rank 3 whose rational minimum rank is 4.

We note that Li et al. [20] showed that for every \(n \times m\) sign pattern with minimum rank 2, there is an integer matrix in its sign pattern class each of whose entries has absolute value at most \(2n - 3\) that achieves the minimum rank.

One goal of this paper is to show that for any \(m \times n\) sign pattern \(A\) with minimum rank \(n - 2\), rational realization of the minimum rank is possible. This is done using a new approach involving sign vectors and duality. Furthermore, it is shown that for each integer \(n \geq 9\), there exists a nonnegative integer \(m\) such that there exists an
Another goal is to use sign vectors of subspaces to investigate the minimum ranks of sign patterns further. In particular, a characterization of \( m \times n \) sign patterns \( A \) with minimum rank \( n - 1 \) is given, which solves an open problem posed in Brualdi et al. [6]. We also obtain a characterization of \( L \)-matrices using sign vectors and a more general description of sign patterns with minimum rank \( r \), in terms of sign vectors of certain subspaces. Several related open problems are also discussed.

2. Sign vectors and duality.

For any vector \( x \in \mathbb{R}^n \), we define the sign vector of \( x \), \( \text{sign}(x) \in \{+, -, 0\}^n \), by

\[
\text{sign}(x)_i = \begin{cases} 
+ & \text{if } x_i > 0, \\
0 & \text{if } x_i = 0, \\
- & \text{if } x_i < 0.
\end{cases}
\]

For any subspace \( L \subseteq \mathbb{R}^n \), we define the set of sign vectors of \( L \) as

\[
\text{sign}(L) = \{ \text{sign}(x) \mid x \in L \}.
\]

Observation 2.1. If \( K \) and \( L \) are subspaces of \( \mathbb{R}^n \) with \( \text{sign}(K) = \text{sign}(L) \), then \( \dim(K) = \dim(L) \).

Indeed, consider the sign vectors of the columns of the reduced column echelon form of a matrix whose columns form a basis of \( K \). Such sign vectors are in \( \text{sign}(K) \) and hence also in \( \text{sign}(L) \). It follows that \( \dim(L) \geq \dim(K) \). By reversing the roles of \( K \) and \( L \), we get the reverse inequality.

For a subspace \( L \subseteq \mathbb{R}^n \), as usual, \( L^\perp = \{ x \in \mathbb{R}^n \mid x^T y = 0 \text{ for all } y \in L \} \) denotes the orthogonal complement of \( L \).

Two sign vectors \( c, x \in \{+, -, 0\}^n \) are said to be orthogonal, written as \( c \perp x \), if one of the following two conditions holds:

1. for each \( i \), we have \( c_i = 0 \) or \( x_i = 0 \), or
2. there are indices \( i, j \) with \( c_i = x_i \neq 0 \) and \( c_j = -x_j \neq 0 \).

For a set of sign vectors \( S \subseteq \{+, -, 0\}^n \), the orthogonal complement of \( S \) is

\[
S^\perp = \{ c \in \{+, -, 0\}^n \mid c \perp x = 0 \text{ for all } x \in S \}.
\]

Notice that if \( c, x \in \mathbb{R}^n \) and \( c^Tx = 0 \), then \( \text{sign}(c) \perp \text{sign}(x) = 0 \).
We will use the following theorem, a proof of which can be found in Ziegler [27].

**Theorem 2.2.** [Duality of oriented matroids] For any subspace $L \subseteq \mathbb{R}^n$,

$$\text{sign}(L)^\perp = \text{sign}(L^\perp).$$

A subspace $L \subseteq \mathbb{R}^n$ is called rational if $L$ has a basis consisting of rational vectors.

**Lemma 2.3.** Let $L$ be a rational subspace of $\mathbb{R}^n$. For any sign vector $s \in \text{sign}(L)$, there exists a rational vector $x \in L$ such that $\text{sign}(x) = s$.

**Proof.** Let $s \in \text{sign}(L)$. Let $B = [b_1 \ b_2 \ \cdots \ b_k]$ be a matrix whose columns form a rational basis for $L$, let $N = \{i : s_i = 0\}$ and let $V = \{z \in \mathbb{R}^k : (Bz)_i = 0 \text{ for all } i \in N\}$. Since $\{b_1, b_2, \ldots, b_k\}$ is a rational basis for $L$, $V$ has a rational basis $\{c_1, c_2, \ldots, c_t\}$. There exists a vector $y \in V$ such that $\text{sign}(By) = s$. Let $a = \begin{bmatrix} a_1 \\ a_2 \\ \cdots \\ a_t \end{bmatrix}^T \in \mathbb{R}^t$ such that $a_1c_1 + a_2c_2 + \cdots + a_tc_t = y$. There exists a sequence of rational vectors $a(n) = \begin{bmatrix} a_1(n) \\ a_2(n) \\ \cdots \\ a_t(n) \end{bmatrix}^T$, $n \in \mathbb{N}$ such that $\lim_{n \to \infty} a(n) = a$. Define $y(n) = a_1(n)c_1 + a_2(n)c_2 + \cdots + a_t(n)c_t$. Then $\lim_{n \to \infty} y(n) = y$ and each $y(n)$ is a rational vector. Since $\lim_{n \to \infty} By(n) = By$, there exists an $m$ such that $\text{sign}(By(m)) = \text{sign}(By) = s$. By letting $x = By(m)$, the statement of the lemma is proved. \[\square\]

**Theorem 2.4.** For all nonnegative integers $m$ and $r$, the following are equivalent:

1. Rational realization of the minimum rank is possible for every $m \times n$ sign pattern matrix with minimum rank $r$.
2. For every subspace $L \subseteq \mathbb{R}^m$ with dimension $r$, there exists a rational subspace $K \subseteq \mathbb{R}^m$ with $\text{sign}(K) = \text{sign}(L)$.

**Proof.** Suppose that rational realization of the minimum rank is possible for every $m \times n$ sign pattern matrices with minimum rank $r$. Let $L \subseteq \mathbb{R}^m$ have dimension $r$. Let $s_1, \ldots, s_k$ be all sign vectors in $\text{sign}(L)$. Let $A$ be the $m \times k$ sign pattern matrix whose $i$th column is $s_i$. Then $\text{mr}(A) \leq r$, as for each sign vector $s_i$ we can choose a vector $x_i \in L$ with $\text{sign}(x_i) = s_i$, and the matrix $B$, whose $i$th column is $x_i$, has rank $r$. Thus, there exists a rational matrix $C \in Q(A)$ with rank at most $r$. Let $K$ be the column space of $C$. Then $\text{sign}(K) = \text{sign}(L)$ and $K$ is rational.

Conversely, suppose that for each subspace $L \subseteq \mathbb{R}^m$ with dimension $r$ there exists a rational subspace $K \subseteq \mathbb{R}^m$ with $\text{sign}(K) = \text{sign}(L)$. Let $A$ be an $m \times n$ sign pattern matrix with $\text{mr}(A) = r$ and let $A_i$ denote the $i$th column of $A$. Since $\text{mr}(A) = r$, there exists a matrix $F \in Q(A)$ with rank($F$) = $r$. Let $L \subseteq \mathbb{R}^m$ be the column space of $F$. Note that $\text{dim}(L) = r$. By assumption, there exists a rational subspace $K$ of $\mathbb{R}^m$ such that $\text{sign}(L) = \text{sign}(K)$. By Observation 2.1, $\text{dim}(K) = r$. Since
sign($L$) = sign($K$), there exist rational vectors $x_i \in K$, $i = 1, 2, \ldots, n$ such that sign($x_i$) = $A_i$, by Lemma 2.3. Let $C = [x_1 \ x_2 \ \cdots \ \ x_n]$. Then $C \in Q(A)$, and $C$ is a rational matrix. Since dim($K$) = $r$ and mr($A$) = $r$, rank $C = r$. □

The next theorem can be found in [11] and [20].

**Theorem 2.5.** Rational realization of the minimum rank is possible for every $m \times n$ sign pattern matrix $A$ with minimum rank 2.

The following theorem follows from the two preceding theorems.

**Theorem 2.6.** For any subspace $L \subseteq \mathbb{R}^n$ with dim($L$) = 2, there exists a rational subspace $K \subseteq \mathbb{R}^n$ with dim($K$) = 2 and sign($K$) = sign($L$).

Since the solution space of a system of homogeneous linear equations with rational coefficients has a rational basis, it is clear that if $L \subseteq \mathbb{R}^n$ is a rational subspace, then $L^\perp$ is also a rational subspace.

**Lemma 2.7.** For any subspace $L \subseteq \mathbb{R}^n$ with dim($L$) = $n - 2$, there exists a rational subspace $K \subseteq \mathbb{R}^n$ with dim($K$) = $n - 2$ such that sign($K$) = sign($L$).

**Proof.** By Theorem 2.6 there exists a rational subspace $M \subseteq \mathbb{R}^n$ with dim($M$) = 2 such that sign($L^\perp$) = sign($M$). Let $K = M^\perp$. Then $K$ is a rational subspace of $\mathbb{R}^n$ with dim($K$) = $n - 2$ and, by Theorem 2.2, sign($K$) = sign($M^\perp$) = sign($M$) = sign($L^\perp$) = sign($L$). □

By Theorem 2.4 and by considering the transpose if needed, we get the following result.

**Theorem 2.8.** Rational realization of the minimum rank is possible for every $m \times n$ sign pattern matrix with minimum rank $m - 2$ or $n - 2$.

From the $9 \times 9$ example given in [18], the next theorem follows immediately.

**Theorem 2.9.** For each integer $n \geq 9$, there exists an $m \times n$ sign pattern matrix $A$ with mr($A$) = 3 for which no rational realization is possible.

From Theorem 2.4, we then obtain the following corollary.

**Corollary 2.10.** Let $n \geq 9$ be an integer. Then there exists a subspace $L \subseteq \mathbb{R}^n$ with dim($L$) = 3 such that there is no rational subspace $K \subseteq \mathbb{R}^n$ with dim($K$) = 3 and sign($L$) = sign($K$).

The following result then follows from Theorem 2.2.

**Lemma 2.11.** Let $n \geq 9$ be an integer. There exists a subspace $M \subseteq \mathbb{R}^n$ with dim($M$) = $n - 3$ such that there is no rational subspace $K \subseteq \mathbb{R}^n$ with sign($K$) =
By Corollary 2.10, there exists a subspace $L \subseteq \mathbb{R}^n$ with $\dim L = 3$ such that there is no rational subspace $K \subseteq \mathbb{R}^n$ with $\dim K = 3$ and $\text{sign}(L) = \text{sign}(K)$. Let $M = L^\perp$. Suppose for a contradiction that there exists a rational subspace $S \subseteq \mathbb{R}^n$ with $\dim(S) = n - 3$ and $\text{sign}(S) = \text{sign}(M)$. Then $S^\perp$ is a rational subspace of $\mathbb{R}^n$ with $\dim(S^\perp) = 3$. By Theorem 2.2, $\text{sign}(S^\perp) = \text{sign}(S)^\perp = \text{sign}(M)^\perp = \text{sign}(M) = \text{sign}(L)$. This contradicts the assumption that there is no rational subspace $K \subseteq \mathbb{R}^n$ with $\text{sign}(L) = \text{sign}(K)$. \(\Box\)

Another application of Theorem 2.4 gives the following fact.

**Theorem 2.12.** For each integer $n \geq 9$, there exists a nonnegative integer $m$ such that there exists an $m \times n$ sign pattern matrix with minimum rank $n - 3$ for which rational realization is not possible.

However, this leaves open the following question.

**Problem 2.13.** Is it true that for each integer $n \geq 9$, there exists an $n \times n$ sign pattern matrix with minimum rank $n - 3$ for which rational realization is not possible?

Another natural question is the following.

**Problem 2.14.** Is it true that for every 3-dimensional subspace $L$ of $\mathbb{R}^n$, there is a rational subspace $K$ such that $\text{sign}(L) = \text{sign}(K)$?

There are connections between rational realization of the minimum ranks of sign patterns and the existence of rational solutions of certain matrix equations, as indicated in [3]. Using this connection and Theorem 2.8 for the case of $\text{mr}(A) = n - 2$, we are able to prove the following result.

**Theorem 2.15.** Suppose that $B$, $C$ and $E$ are real matrices such that $BC = E$. If $E$ has either 2 rows or 2 columns, then there exist rational matrices $\tilde{B}$, $\tilde{C}$ and $\tilde{E}$ such that $\text{sign}(\tilde{B}) = \text{sign}(B)$, $\text{sign}(\tilde{C}) = \text{sign}(C)$, $\text{sign}(\tilde{E}) = \text{sign}(E)$, and $\tilde{B}\tilde{C} = \tilde{E}$.

*Proof.* Without loss of generality, assume that $E$ has two columns. Consider the $2 \times 2$ block matrix

$$M = \begin{bmatrix} I_n & C \\ B & E \end{bmatrix}. $$

Observe that $M$ has $n + 2$ columns and the Schur complement of $I_n$ in $M$ is 0. Hence, $\text{rank}(M) = n$. It follows that the minimum rank of the sign pattern $\text{sign}(M)$ is $n$. Hence, from Theorem 2.8 there is a rational matrix

$$\tilde{M} = \begin{bmatrix} D & C_1 \\ B_1 & E_1 \end{bmatrix}.$$
of rank \( n \) in \( Q(\text{sign}(M)) \).

It follows that the Schur complement of \( D \) in \( \tilde{M} \) is \( E_1 - B_1D^{-1}C_1 = 0 \). The rational matrices \( \tilde{B} = B_1, \ \tilde{C} = D^{-1}C_1 \) and \( \tilde{E} = E_1 \) clearly satisfy the desired properties. \( \square \)

Since the zero entries of the matrix \( E \) in a real matrix equation \( BC = E \) are the major obstructions for the existence of rational solutions within the same corresponding sign pattern classes, the preceding theorem suggests the following conjecture.

**Conjecture 2.16.** Let \( B, C \) and \( E \) be real matrices such that \( BC = E \). If all the zero entries of \( E \) are contained in a submatrix with either 2 rows or 2 columns, then there exist rational matrices \( \tilde{B}, \tilde{C}, \) and \( \tilde{E} \), such that \( \text{sign}(\tilde{B}) = \text{sign}(B), \ \text{sign}(\tilde{C}) = \text{sign}(C), \ \text{sign}(\tilde{E}) = \text{sign}(E), \) and \( \tilde{B}\tilde{C} = \tilde{E} \).

### 3. Further results on minimum ranks and sign vectors of subspaces.

We now investigate the minimum ranks of sign patterns further using sign vectors and duality.

Similar to the concept of \( \text{mr}(A) \), the **maximum rank** of a sign pattern matrix \( A \), denoted \( \text{MR}(A) \), is the maximum of the ranks of the real matrices in \( Q(A) \). It is well-known that \( \text{MR}(A) \) is the maximum number of nonzero entries of \( A \) with no two of the nonzero entries in the same row or in the same column. By a theorem of König, the minimal number of lines (namely, rows and columns) in \( A \) that cover all of the nonzero entries of \( A \) is equal to the maximal number of nonzero entries in \( A \), no two of which are on the same line. This common number is also called the **term rank** \( [7, 15] \). In contrast to the rational realization problem of the minimum rank, we note that through diagonal dominance, it can be easily seen that for every sign pattern matrix \( A \), \( \text{MR}(A) \) can always be achieved by a rational matrix.

Sign pattern matrices \( A \) that require a unique rank (namely, \( \text{MR}(A) = \text{mr}(A) \)) were characterized by Hershkowitz and Schneider in \([16]\).

It is shown in \([2]\) that rational realization of the minimum rank of \( A \) is always possible if \( \text{MR}(A) - \text{mr}(A) = 1 \).

In \([2]\), the study of sign patterns \( A \) with \( \text{MR}(A) - \text{mr}(A) = 1 \) is reduced to the study of \( m \times n \) sign patterns \( A \) such that \( \text{MR}(A) = n \) and \( \text{mr}(A) = n - 1 \). In this connection, using sign vectors and duality, we obtain the following characterization of the \( m \times n \) sign patterns \( A \) such that \( \text{MR}(A) = n \) and \( \text{mr}(A) = n - 1 \) (which was raised as an open problem in \([5]\)). Since the condition \( \text{MR}(A) = n \) is easily checked, it suffices to characterize \( m \times n \) sign patterns \( A \) with \( \text{mr}(A) = n - 1 \).

For an \( m \times n \) sign pattern or real matrix, we denote by \( \text{row}(A) \) the set of the row
vectors of $A$.

**Theorem 3.1.** Let $A$ be an $m \times n$ sign pattern. Then $\mr(A) = n - 1$ if and only if the following two conditions hold:

(a) There is a nonzero sign vector $x \in \{+,-,0\}^n$ such that $\row(A) \subseteq x^\perp$.

(b) For every two-dimensional subspace $L$ of $\mathbb{R}^n$, $\row(A) \not\subseteq \sign(L)^\perp$.

Proof. Assume that $\mr(A) = n - 1$. Then there is a real matrix $A_1 \in Q(A)$ such that $\rank(A_1) = n - 1$. Let $v_0$ be a nonzero vector in the null space $\nullspace(A_1)$ and let $x = \sign(v_0)$. Then it is clear that $\row(A) \subseteq x^\perp$, namely, (a) holds. Further, consider any two-dimensional subspace $L$ of $\mathbb{R}^n$. Suppose that $\row(A) \subseteq \sign(L)^\perp$. Then by Theorem 2.2, $\row(A) \subseteq \sign(L^\perp)$. Since $L^\perp$ is a subspace of dimension $n - 2$ and every row vector of $A$ is contained in $\sign(L^\perp)$, it is clear that the span of the real vectors in $L^\perp$ whose sign vectors agree with the rows of $A$ is a subspace of dimension at most $n - 2$, that is to say, there is a real matrix $A_2 \in Q(A)$ of rank at most $n - 2$, which contradicts the assumption that $\mr(A) = n - 1$.

Conversely, assume that (a) and (b) hold. Let $x_1$ be the $(1,-1,0)$ vector in $Q(x)$. From (a), it is easily seen that there is a real matrix $A_1 \in Q(A)$ such that $A_1 x_1 = 0$. It follows that $\rank(A_1) \leq n - 1$ and hence $\mr(A) \leq n - 1$. Suppose that $\mr(A) \leq n - 2$. Then there is a matrix $A_2 \in Q(A)$ with $\rank(A_2) \leq n - 2$. It follows that $\nullspace(A_2)$ has dimension at least 2. Hence, $\nullspace(A_2)$ contains a subspace $L \in \mathbb{R}^n$ of dimension 2. Since $L \subseteq \nullspace(A_2)$, we have that $\row(A_2) \subseteq L^\perp$, and hence $\row(A) = \sign(\row(A_2)) \subseteq \sign(L^\perp)$. By Theorem 2.2, $\sign(L^\perp) = \sign(L)^\perp$, and so $\row(A) \subseteq \sign(L)^\perp$, contradicting (b). Thus, $\mr(A) = n - 1$. $
$

We note that for an $m \times n$ sign pattern matrix $A$, the condition that each row vector of $A$ is in $\{u,v\}^\perp$ for some two nonzero sign vectors $u, v \in \{+,-,0\}^n$ with $u \neq \pm v$ does not imply that $\mr(A) \leq n - 2$, as the following example shows.

**Example 3.2.** Let $A = \begin{bmatrix} + & + & + \\ 0 & + & + \end{bmatrix}$, $u = \begin{bmatrix} + \\ + \\ 0 \\ - \end{bmatrix}$, and $v = \begin{bmatrix} 0 \\ + \end{bmatrix}$. It is obvious that each row vector of $A$ is in $\{u,v\}^\perp$ for the two nonzero sign vectors $u, v \in \{+,-,0\}^3$ with $u \neq \pm v$. But clearly, $\mr(A) = 2$, so $\mr(A) \nleq n - 2 = 1$.

Using a similar argument as in the proof of Theorem 3.1, we can show the following more general result that characterizes sign patterns with minimum rank $r$, for each possible $r \geq 1$.

**Theorem 3.3.** Let $A$ be an $m \times n$ sign pattern and let $\row(A)$ denote the set of the row vectors of $A$, and let $r$ be any integer such that $1 \leq r \leq \min\{m,n\}$. Then $\mr(A) = r$ if and only if the following two conditions hold:
There is a subspace $L \subseteq \mathbb{R}^n$ with $\dim(L) = r$ such that $\text{row}(A) \subseteq \text{sign}(L)$.

(b) For every subspace $V$ of $\mathbb{R}^n$ with $\dim(V) = r - 1$, $\text{row}(A) \not\subseteq \text{sign}(V)$.

In particular, the above theorem gives a characterization of $L$-matrices (namely, $m \times n$ sign patterns $A$ with $\text{mr}(A) = n$). Note that every subset of $\{+,-,0\}^n$ is of course contained in $\text{sign}(\mathbb{R}^n)$ and Theorem 3.5 gives us that for every subspace $V$ of $\mathbb{R}^n$ of dimension $n - 1$, $\text{sign}(V) = \text{sign}(x^\perp)$ for each nonzero vector $x$ in $V^\perp$.

**Corollary 3.4.** Let $A$ be an $m \times n$ sign pattern and let $\text{row}(A)$ denote the set of the rows of $A$. Then, $\text{mr}(A) = n$ if and only if for every nonzero sign vector $x \in \{+,-,0\}^n$, $\text{row}(A) \not\subseteq x^\perp$.

The last characterization of $L$-matrices can be seen to be equivalent to the characterization of such matrices found in [8] (where a sign pattern is defined to be an $L$-matrix if the minimum rank is equal to its number of rows).

Let $S_{k,n}$ (respectively, $s_{k,n}$) denote the maximum cardinality (respectively, minimum cardinality) of $\text{sign}(L)$ as $L$ runs over all $k$-dimensional subspaces of $\mathbb{R}^n$. In other words,

$$S_{k,n} = \max_{L \subseteq \mathbb{R}^n, \dim(L) = k} |\text{sign}(L)|, \quad s_{k,n} = \min_{L \subseteq \mathbb{R}^n, \dim(L) = k} |\text{sign}(L)|.$$

For every subspace $L$, as the nonzero vectors in $\text{sign}(L)$ occur in disjoint pairs of vectors that are negatives of each other, it is clear that $|\text{sign}(L)|$ is odd. Thus, $S_{k,n}$ and $s_{k,n}$ are always odd. For each $k$ ($0 \leq k \leq n - 1$), since every $k$-dimensional subspace of $\mathbb{R}^n$ is contained in a subspace of dimension $k + 1$, it is clear that $S_{k,n} \leq S_{k+1,n}$ and $s_{k,n} \leq s_{k+1,n}$.

Obviously, $S_{0,n} = s_{0,n} = 1$, $S_{1,n} = s_{1,n} = 3$, and $S_{n,n} = s_{n,n} = 3^n$. By considering the reduced row echelon form of a matrix $A$ whose rows form a basis for a $k$-dimensional subspace $L$ of $\mathbb{R}^n$ and observing that the components of the vectors in $L$ in the pivot columns of $A$ are independent and arbitrary, it can be seen that $s_{k,n} \geq 3^k$, for each $k$, $1 \leq k \leq n$. For the $k$-dimensional subspace $L$ spanned by the standard vectors $e_1, e_2, \ldots, e_k$, it can be seen that equality in the last inequality can be achieved. Thus, $s_{k,n} = 3^k$, for each $k$, $1 \leq k \leq n$. We record this result below.

**Theorem 3.5.** Let $n \geq 2$. Then $s_{k,n} = 3^k$, for each $k$, $1 \leq k \leq n$.

Let $L$ be a $k$-dimensional linear subspace of $\mathbb{R}^n$. Let $B$ be a matrix whose columns form a basis of $L$. For each $1 \leq j \leq n$, let $H_j$ be the orthogonal complement of the $j$th row of $B$ in $\mathbb{R}^k$. The central hyperplane arrangement $\{H_1, \ldots, H_n\}$ in $\mathbb{R}^k$ partitions $\mathbb{R}^k$ into disjoint, relatively open cells of dimensions 0 through $k$, with each cell corresponding to precisely one sign vector in $\text{sign}(L)$, see [21, 23, 25, 27]. It can be seen that $S_{k,n}$ is equal to the total number of cells of a generic central hyperplane...
arrangement \( \{ H_1, \ldots, H_n \} \) in \( \mathbb{R}^k \). For \( n \geq k \), adding a hyperplane \( H_{n+1} \) to a generic central hyperplane arrangement \( \{ H_1, \ldots, H_n \} \) in generic position increases the total number of cells by \( 2(S_{k-1,n} - 1) \). This can be seen as follows. For \( m > 1 \), each relatively open cell of dimension \( m \) of the central hyperplane arrangement \( \{ H_1, \ldots, H_n \} \) that is intersected by \( H_{n+1} \) yields three relatively open cells in the hyperplane arrangement \( \{ H_1, \ldots, H_{n+1} \} \), one of dimension \( m - 1 \), which is a relatively open cell of \( \{ H_1 \cap H_{n+1}, \ldots, H_n \cap H_{n+1} \} \), and two of dimension \( m \). For \( m \leq 1 \), each relatively open cell of dimension \( m \) of the central hyperplane arrangement \( \{ H_1, \ldots, H_n \} \) is not cut by \( H_{n+1} \) and is a relatively open cell in the hyperplane arrangement \( \{ H_1, \ldots, H_{n+1} \} \).

(Here, we used that \( n \geq k \), so that there exists a cell of dimension 0 of the arrangement \( \{ H_1, \ldots, H_n \} \).) Since the number of relatively open cells of the arrangement \( \{ H_1 \cap H_{n+1}, \ldots, H_n \cap H_{n+1} \} \) in \( H_{n+1} \) is \( S_{k-1,n} \), the increase of the total number of cells is \( 2(S_{k-1,n} - 1) \). Therefore the following recursion formula holds for each \( n \geq 2 \) and each \( 1 \leq k < n \):

\[
S_{k,n} = S_{k,n-1} + 2(S_{k-1,n-1} - 1).
\]

As pointed out by Richard Stanley, the total number of cells of a generic central hyperplane arrangement \( \{ H_1, \ldots, H_n \} \) of \( \mathbb{R}^k \) can be computed using the intersection lattice approach described in [24, 26]. The resulting formula is given below.

**Theorem 3.6.** For each \( n \geq 2 \) and each \( 1 \leq k \leq n \),

\[
S_{k,n} = 1 + \sum_{i=0}^{k-1} 2^i \binom{n}{i} + \sum_{i=1}^{k} \binom{n}{k-i} \binom{n-k+i-1}{i-1}.
\]

**Example 3.7.**

1. \( S_{1,n} = 3 \).
2. \( S_{2,n} = 4n + 1 \).
3. \( S_{3,n} = 4n^2 - 4n + 3 \).
4. \( S_{4,n} = \frac{8}{3}n^3 - 8n^2 + \frac{28}{3}n + 1 \).
5. \( S_{5,n} = \frac{4}{3}n^4 - 8n^3 + \frac{64}{3}n^2 - 12n + 3 \).
6. \( S_{6,n} = \frac{32}{15}n^5 - \frac{16}{3}n^4 + \frac{64}{3}n^3 - \frac{404}{3}n^2 + \frac{332}{15}n + 1 \).

Consequently, we have the following results on sign pattern matrices.

**Theorem 3.8.** Let \( A \) be an \( m \times n \) sign pattern matrix.

1. If \( mr(A) = 1 \), then \( |row(A)| \leq 3 \).
2. If \( mr(A) = 2 \), then \( |row(A)| \leq 4n + 1 \).
3. If \( mr(A) = 3 \), then \( |row(A)| \leq 4n^2 - 4n + 3 \).
4. If \( mr(A) = 4 \), then \( |row(A)| \leq \frac{8}{3}n^3 - 8n^2 + \frac{28}{3}n + 1 \).
5. If \( mr(A) = 5 \), then \(|\text{row}(A)| \leq \frac{4}{15}n^5 - \frac{16}{7}n^4 + \frac{64}{3}n^3 - \frac{404}{33}n^2 + \frac{332}{3}n + 1. \)

6. If \( mr(A) = 6 \), then \(|\text{row}(A)| \leq \frac{8}{75}n^5 - \frac{16}{7}n^4 + \frac{64}{3}n^3 - \frac{404}{33}n^2 + \frac{332}{3}n + 1. \)

**Final remarks.** The main results of this paper were contained in a preliminary version of the paper, and were presented at the 2013 ILAS conferences by one of the coauthors. In 2014, we noticed the related paper by Shitov [23] that has some overlap. Unfortunately, the proof of a key lemma (Lemma 3.5) in Shitov’s paper has a logical gap.

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**REFERENCES**


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